

Cyclic Redundancy Check (CRC) Example, Part 2

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1. An Example with CRC-CCITT

The CRC-CCITT generator polynomial is

$$\text{In[289]:= } g = x^{16} + x^{12} + x^5 + 1$$

$$\text{Out[289]= } 1 + x^5 + x^{12} + x^{16}$$

Let's assume we are encoding the message bit stream 0005040010₁₆. The binary digits will be the coefficients of our message polynomial $i(x)$, with the least significant bit being the constant term of the polynomial:

$$\text{In[290]:= } i = x^{26} + x^{24} + x^{18} + x^4 + x$$

$$\text{Out[290]= } x + x^4 + x^{18} + x^{24} + x^{26}$$

Enter the code's blocklength n and message length k , and compute the the number of parity bits $n-k$.

$$\text{In[291]:= } n = 2^{15} - 1$$

$$\text{Out[291]= } 32767$$

$$\text{In[292]:= } k = n - 16$$

$$\text{Out[292]= } 32751$$

$$\text{In[293]:= } n - k$$

$$\text{Out[293]= } 16$$

For systematic encoding, the parity is $p(x) = [-x^{n-k} i(x)] \bmod g(x)$ where we do modulo 2 arithmetic on the polynomial coefficients.

$$\text{In[294]:= } p = \text{PolynomialMod}[x^{n-k} i, \{g, 2\}]$$

$$\text{Out[294]= } x + x^2 + x^4 + x^5 + x^6 + x^7 + x^9 + x^{12} + x^{14}$$

Writing the polynomial coefficients in binary we get the parity,
0101 0011 1111 0110 _{2=53F5} ₁₆

Compute the systematically encoded codeword

$$c(x) = x^{n-k} i(x) + p(x).$$

$$\text{In[295]:= } c = \text{Expand}[x^{n-k} i + p]$$

$$\text{Out[295]= } x + x^2 + x^4 + x^5 + x^6 + x^7 + x^9 + x^{12} + x^{14} + x^{17} + x^{20} + x^{34} + x^{40} + x^{42}$$

Compute the shifted syndrome $s'(x) = [x^{n-k} c(x)] \bmod g(x)$, modulo 2 on the polynomial coefficients. We should get zero.

```
In[296]:= p = PolynomialMod[xn-k c, {g, 2}]
```

```
Out[296]= 0
```

Add error to the codeword.

```
In[297]:= cerror1 = c + x11
```

```
Out[297]= x + x2 + x4 + x5 + x6 + x7 + x9 + x11 + x12 + x14 + x17 + x20 + x34 + x40 + x42
```

Compute the shifted syndrome $s'(x) = [x^{n-k} c(x)] \bmod g(x)$ modulo 2 on the polynomial coefficients. We should get a non-zero answer.

```
In[298]:= s = PolynomialMod[xn-k cerror1, {g, 2}]
```

```
Out[298]= 1 + x3 + x5 + x7 + x8 + x11 + x15
```

On the other hand, if we add a multiple of a codeword, we won't see the error.

```
In[299]:= cerror2 = c + x2 g
```

```
Out[299]= x + x2 + x4 + x5 + x6 + x7 + x9 + x12 + x14 + x17 + x20 + x34 + x40 + x42 + x2 (1 + x5 + x12 + x16)
```

```
In[300]:= s = PolynomialMod[xn-k cerror2, {g, 2}]
```

```
Out[300]= 0
```